

Parallelism 2

Note Title

12/10/2008

Return marked up prog. 3 to envelope on by door by 12/19
5 PM

Performance test case on website tonight.

Course Evals

Multicore is wave of future.

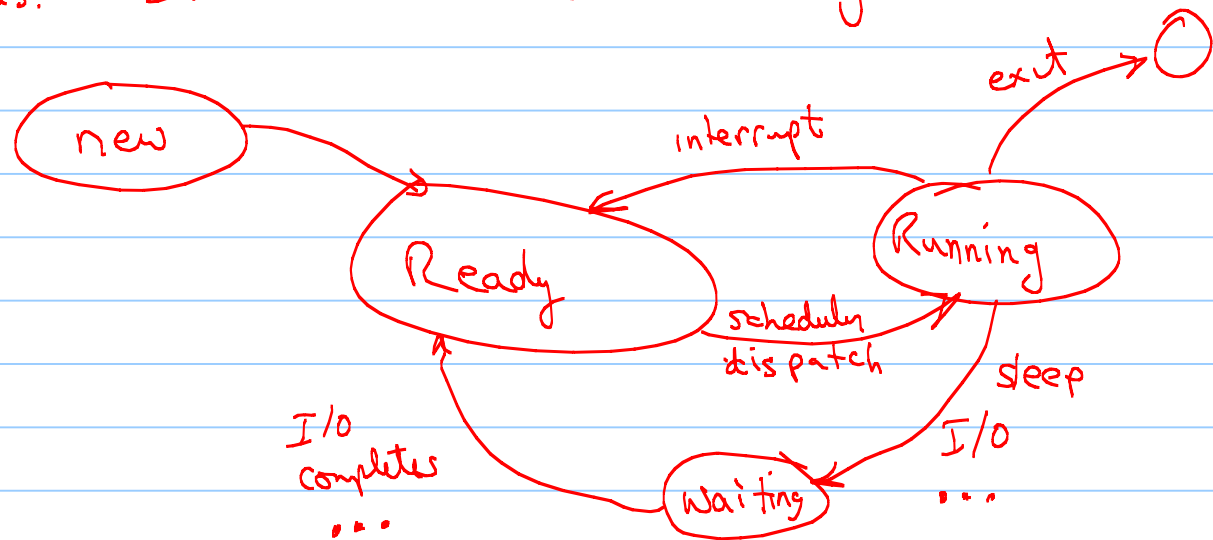
Multi. threading

- Goal is to make a single app use multiple cores
- A Thread (of execution) is a sequential computation w/ its own PC, stack, registers, but not its own memory address space.

Contrast: a process is the same but w/ own A.S.

Idea: each thread can run on a different core.

Many more threads than cores. How does O.S. manage threads. State Machine for scheduling



Threads share address space — can access same locations.

T1

read x
inc x (in ^{private} reg)
write x

T2

read x
inc x (in private reg)
write x

x: 10

Race - condition.

Programming w/ threads have to use synchronization
to prevent race conditions.

A lock (mutual exclusion device, mutex) is a guard on a critical section.

T1

lock(x)

read x

inc x

write x

unlock(x)

T2

=

=

=

=

=

When T2 tries to lock x and T1 has already locked x, T2 goes to wait state.

T1
lock(x)
:
x
lock y
:
x and y
unlock(x, y)

T2
lock(y)
:
y
lock x
:
x and y
unlock(x, y)

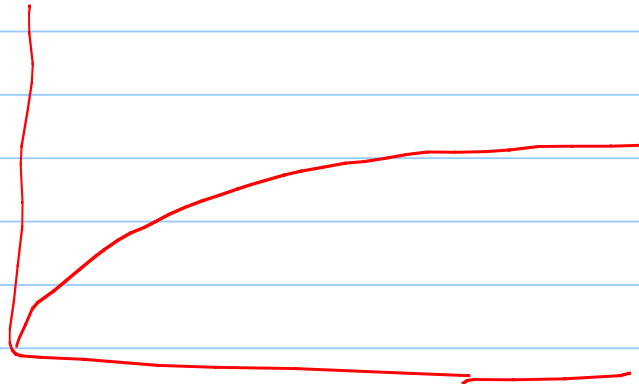
Dead lock !

End result: programing w/ threads is not at all
like programing sequentially



Gene Amdahl :

Amdahl's law



Every program has some sequential fraction = S
Speedup obtainable w/ P processors =

$$\text{Time}(1 \text{ proc}) / \text{Time}(P \text{ procs}) \leq \frac{1}{S + ((1-S)/P)} \text{ as } P \rightarrow \infty : 1/S$$